What is claimed is:

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1	1.	A method for use in a database system, comprising:
2 /		storing a materialized join view based on at least two base
3	relations;	
4		storing at least one auxiliary relation containing one or more
5	attributes of	one of the base relations, the auxiliary relation partitioned according
6	to a join attr	ibute; and
7		updating the at least one auxiliary relation in response to
8	modification	of the one base relation.
1	2.	The method of claim 1, further comprising storing the one base
2	relation that	is not partitioned according to the join attribute.
1	3.	The method of claim 2, further comprising:
2		receiving a tuple into the database system;
3		storing the tuple in one of the at least two base relations;
4		storing the tuple in the auxiliary relation; and
5		using the auxiliary relation to determine whether to update the
6	materialized	join view.
1	4.	A method comprising:
2	•	receiving a tuple into a relation at a first node, wherein the tuple
3	comprises a	join attribute and the relation is not partitioned according to the join
4	attribute;	
5		storing the tuple in an auxiliary relation at a second node, wherein
6	the auxiliary	relation is partitioned according to the join attribute;
7		identifying second tuples of a second relation;
8		joining the tuple with the second tuples to produce join results; and
9		storing the join results in a join view.

1	5.	The method of claim 4, wherein joining the tuple with the second
2	tuples to pro	oduce join results comprises:
3		identifying second join attributes in the second tuples, and
4		comparing the second join attributes with the join attribute of the
5	relation.	
1	6.	The method of claim 5, wherein identifying second tuples of a
2	second relat	ion comprises identifying second tuples of the second relation at the
3	second node	e. /
1	7.	The method of claim 4, wherein storing the tuple in an auxiliary
2	relation at a	second node comprises:
3		determining that a join view definition excludes an attribute of the
4	tuple; and	
5		not storing the excluded attribute in the auxiliary relation.
1	8.	The method of claim 4, wherein storing the tuple in an auxiliary
2	relation at a	second node comprises: /
3		determining that a join/view definition includes a condition on one
4	of the attribu	utes of the tuple; and /
5		identifying the attribute in the tuple;
6		determining that the condition is not met; and
7		not storing the tuple in the auxiliary relation.
1	9.	The method of claim 4, wherein storing the tuple in an auxiliary
2	relation at a	second node comprises:
3		determining that a join view definition includes a condition that
4	cannot be m	et; and /
5		not storing the tuple in the auxiliary relation.

1	10.	The method of claim 5, wherein storing the tuple in an auxiliary
2	relation at a	second node comprises:
3		determining that the join attribute is a key of the relation;
4		determining that the second join attribute is a foreign key of a
5	second relat	tion, wherein the foreign key references to the join attribute; and
6		not maintaining a second auxiliary relation.
1	11.	The method of claim 4, wherein storing the join results in a join
2	view compri	ises:
3		determining that a join view definition includes a condition on one
4	of the attrib	utes of the tuple; and
5		identifying the attribute in the tuple;
6		determining that the condition is not met; and
7		not storing the join results in the join view.
1	12.	The method of claim 4, wherein storing the join results in a join
2	view compri	ises:
3		determining that a join view definition includes a condition that
4	cann	ot be met; and
5		not storing the join results in the join view.
1	13.	A database system comprising:
2		a storage module to store relations and at least one auxiliary
3	relation cor	responding to one of the relations, the at least one auxiliary relation
4	_	one or more attributes of the one relation, the at least one auxiliary
5	relation par	titioned differently than the one relation, the storage module further
6	to store a jo	oin view based on a/join of the relations; and
7		a controller adapted to update the join view using the at least one
8	auxiliary rel	ation /

			,
	1	14.	The database system of claim 13, wherein the controller is further
	2	adapted to r	receive a tuple and store the tuple in one of the relations and in one
	3	of the at leas	st one auxiliary relations.
	1	15.	The database system of claim 14, wherein the controller is further
	2	adapted to n	not update the join view after receiving some tuples.
	1	16.	An article comprising a medium storing instructions for enabling a
	2	processor-ba	esed system to:
	3		receive a tuple into a relation/at a first node, wherein the tuple
illus i	4	comprises a	join attribute and the relation is not partitioned according to the join
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	5	attribute;	
الساة المساة المساة المساة الساة فلساء فلساة المساة	6		store the tuple in an auxiliary relation at a second node, wherein
	7	the auxiliary	relation is partitioned according to the join attribute;
	8		identify second tuples of a second relation;
1	9		join the tuple with the second tuples to produce join results; and
the that the third that the	10		store the join results/in a join view.
Ŧ	1	17.	The article of claim 16, further storing instructions for enabling the
	2	processor-ba	ased system to:
	3		identify second/join attributes in the second tuples; and
	4		compare the second join attributes with the join attribute of the
	5	relation.	
	1	18.	The article of claim 16, further storing instructions for enabling the
	2	processor-ba	ased system to:
	3		determine that a join view definition excludes an attribute of the
	4	tuple;	
	5		not/store the excluded attribute in the auxiliary relation.
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1	19.	The article of claim 16, further storing instructions for enabling the
2	processor-b	ased system to:
3		determine that a join view definition includes a condition on one of
4	the attribute	es of the tuple; and
5		identify the attribute in the tuple;
6		determine that the condition is not met; and
7		not store the tuple in the auxiliary relation.
1	20.	The article of claim 16, further storing instructions for enabling the
2	processor-b	ased system to:
3		determine that a join view definition includes a condition that
4	cannot be m	net; and
5		not store the tuple in the auxiliary relation.
1	21.	The article of claim 17, further storing instructions for enabling the
2	processor-b	ased system to:
3		determine that the join attribute is a key of the relation;
4		determine that the second join attribute is a foreign key of a
5	second relat	tion, wherein the foreign key references to the join attribute; and
6		not maintain a second auxiliary relation.
1	22.	The article of claim 16, further storing instructions for enabling the
2	processor-b	ased system to: /
3		determine that a join view definition includes a condition on one of
4	the attribute	es of the tuple and
5		identify the attribute in the tuple;
6		determine that the condition is not met; and
7		not store the join results in the join view.

1	23. The article of claim 16, further storing instructions for enabling the
2	processor-based system to:
3	determine that a join view definition includes a condition that
4	cannot be met; and
5	not store the join results in the join view.